Accelerated Simulation Performance through High Performance Computing for Advanced Sealing Applications

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Abstract: The automotive and heavy-duty industries are heavy users of Computer Aided Engineering (CAE) for development, design and performance optimization of their products. As a technology driven company, the Sealing Products Group of Dana Holding Corporation utilizes this technology extensively, especially for metal cylinder head gaskets (CHG) and exhaust gaskets, so called Multi-Layer-Steel (MLS) gaskets. Due to the widespread involvement of complex contact interactions in high fidelity sealing system analysis, high performance computing is essential to accelerating the design process, thus meeting aggressive time and cost to market requirements.

Since the release of Abaqus version 6.6 a major step forward was made through the extensive and effective use of parallel computing. The combination of flexible computing clusters and the power of parallel processing enabled by SIMULIA resulted in significant time and cost savings for Dana when using the Abaqus engineering software package. It also allowed for the implementation of new simulation techniques for contact simulations. This allowed for better and more accurate up-front prediction of sealing technology.

This paper provides an overview on accelerated simulation performance studies utilizing the latest SIMULIA technology for advanced sealing application from the Sealing Products Group of Dana Holding Corp. Furthermore it provides an outlook of steps to come.

Keywords: CAE, High performance computing, cluster, SIMULIA, Abaqus, Dana, sealing products, contact
1. **Introduction**

The requirements for better and more effective performance in development, in particular in the automotive industry became especially important in the current market situation. An ineffective business will not be able to survive the drastic market challenges regarding cost, quality and performance. Computer Aided Engineering (CAE) plays a major role in achieving this effectively. It has been demonstrated that companies focusing heavily on CAE in a coordinated manner are facing less constraints and have a better chance of coming out of this economic down-turn with greater ease and an increased focus than companies that do not utilize CAE.

Within this paper we focus on powertrain applications from a sealing products perspective, as used in the automotive and heavy-duty industries. It outlines the possibilities and different ways to approach CAE in a cost-effective manner driven by high-performance.

1.1 **CAE for Sealing Applications – A brief introduction**

1.1.1 **CAE applications for the sealing of powertrain and exhaust applications**

When talking about CAE applications for the sealing of powertrain and exhaust applications one first needs to know what is understood to be sealing products. Those products include flat gaskets and ring seals, rubber-molded gaskets, seals, covers, pans (thermo-plastic and the thermo-set), and shielding components.

- Flat gaskets and ring seal products, by example:
  - Cylinder head gaskets
  - Rocker cover gaskets
  - Cylinder block hand hole gaskets
  - Exhaust manifold gaskets
  - Turbo charger gaskets
  - Exhaust gas recirculation (EGR) gaskets
  - Slip joint gaskets
  - V-band gaskets

- Rubber-molded gaskets and seal products, by example:
  - Valve cover gaskets
  - Cam cover gaskets
  - Front cover gaskets
  - Window gaskets
- Intake gaskets
- Valve stem seals
- Rubber inserts

• Covers and pan products, by example:
  - Valve covers
  - Cam covers
  - Front covers
  - Oil pans

• Shielding component products, by example:
  - Exhaust manifold shields
  - Down pipe shields
  - Sensor shields

Figure 1 provides a visual example of involved powertrain sealing products.
Those applications range from:

- Low temperatures (below 400°C) to high temperature (above 1000°C)
- Quasi-static to dynamic
- Combustion gas to fluid seals
- Thermal to noise shielding
- Thermal reflection to thermal insulation
- Friction withstanding to friction reducing
- Sealing to controlled metering of fluids

This variety of applications demonstrates the challenges and complexity of CAE for those products including:

- Forming
- Molding
- Contact
- Structural and fluid dynamics simulation and their interaction
- NVH
- Fatigue.

1.1.2 General challenges that need to be considered and addressed

When talking about simulation performance optimization for highly engineered systems we need to understand the different factors of influence. A major measure for this is the so called “sweet spot.” This point represents the optimization between scalability (speed up of a simulation) and cost per unit. Ideally the speed-up factor is increasing linearly with the addition of cores to the simulation. In reality, however, the software has to be capable of utilizing the hardware architecture effectively. This shows the need to benchmark each new hardware and software configuration in order to identify this “sweet spot.”

The diversity of models and their many variations also have an impact on performance. The knowledge of their influence in conjunction with the software and hardware configuration will also have a major impact on the cost and time structure in an “analysis driven” business.
1.1.3 Simulation requirements for sealing applications

The analysis of sealing products for a powertrain and related components involves multidiscipline CAE groups who work together for cooperative analysis including:

- Structural and thermal Stress
- Durability
- Tribology
- CFD

Coupled analysis between CFD and finite element analysis is applied to a wide range of applications in:

- Base engine:
  - In-cylinder
  - Coolant systems
  - Intake manifolds
- Other powertrain components:
  - Air intake systems
  - EGR
  - Turbo systems

Usually the different analysis types are handled within a cross-functional team where there is a large amount of data exchange and reuse such as with material properties, thermal profiles and large meshed models. Challenges with current CAE systems not only include handling the large computational power, but also multiple users and storage considerations for the large amount of data that is being generated.

As more and more analysis is done at a system level rather than at component level or subsystem level, the models are becoming much larger in size and very complex. At the same time the number of “degrees of freedom” (DOF) is increasing to be able to improve the accuracy and response extraction. On top of this, with a reduction in time to market and “first time right” initiatives a much more detailed analysis combined with faster turn-around times is mandated. The types of sealing analysis models can range anywhere from a base engine powertrain assembly model with all of the secondary components included, to a small joint analysis of a subsystem. Depending on the model requirements and also the level of detail required for the results, today’s models sizes may vary from 500,000 DOF to 10 million DOF.
1.2 Simulation performance criteria

In the context of High-Performance Computing (HPC) for finite element analysis (FEA), today, compute clusters are the most common choice. In the last decade, starting with academic and government research labs, compute clusters started to take over the traditional supercomputers and large Shared Memory Multi-Processing (SMP) computers. The main driving force behind this change was the low cost of these clusters which initially were mostly built out of workstation components and running the open source Linux operating system (OS). Compute clusters are characterized as a collection of independent compute nodes, usually with their own local disk and memory subsystem, connected by a high speed interconnect and are often referred to as Distributed Memory Processing (DMP) system. Today customized turn-key compute clusters are available from many hardware vendors and are becoming increasingly popular with the simulation community. While traditionally compute clusters have been more heavily utilized for Computational Fluid Dynamics (CFD) simulations, owing to increasing model sizes and fidelity, they are now also being used for FEA. However, efficient utilization of compute clusters requires Distributed Memory Parallel execution capabilities as well as moreover parallel scalability from any software that seeks to provide a feasible solution. It is therefore important to be able to investigate the feasibility of running FEA simulations on compute clusters for particular workflows.

To calculate the cost of simulations on compute clusters three major contributors must be taken in to account:

- Software
- Hardware
- Engineering time.

Software cost is relatively straightforward to estimate as long as proper utilization and parallel computing licensing costs are employed. On the other hand, calculating hardware cost may be a little more cumbersome. For in-house clusters, the expected life of the cluster, utilization, maintenance, administration, computer room space, power consumption, and cooling costs should all be taken in to account. One emerging alternative to in-house cluster computing is using utility computing services. These services typically charge a fixed amount for each compute node hour used by the customer with prices depending on the precise configuration of the compute cluster offered. Once software and hardware cost per hour are known, it is simply a manner of running representative benchmarks on a range of compute nodes or cores to generate cost curves and determine the cheapest configuration - alternatively referred to as the “sweet spot” - to run simulations. For the purposes of this paper we present a feasibility study using three different simulations and limit ourselves to analyzing the cost using utility computing services from R Systems NA, Inc.
1.3 General Simulation Structure

The recent industry push towards DMP computing has led to the advanced development and scalability of Applied Engineering applications such as Abaqus. This trend has allowed industry engineers to greatly increase simulation performance while reducing the previously high costs associated with large SMP compute systems. Replacing the specialized SMP hardware with the more cost effective, modular, and readily available “node” system DMP models allows for more than just up-front cost savings. In addition to the bottom line price, the modular architecture of today’s compute note DMP systems allows for an almost endless array of hardware configurations and expandability, enabling today’s hardware to meet the flexibility and scalability of the rapidly changing CAE software.

While cluster computing capabilities of Abaqus/Standard were first introduced with release 6.6, significant improvements were delivered with later releases. Just before Abaqus 6.7 was released in 2007, SIMULIA and Dana engaged in the first cluster computing feasibility study. This study demonstrated that to minimize the simulation cost (i.e. hit the sweet spot) cluster computing with at least eight compute nodes was a must. This motivated Dana to quickly adopt an in-house cluster computing solution for their FEA simulations. Such collaborations still continue to date. Most recently, SIMULIA, Dana, and R Systems worked together on a HPC study which resulted in this press release [1] on the performance of state-of-the-art cluster computing platforms.

The current analysis environment is composed of a single high performance DMP compute cluster. Taking advantage of the benefits of DMP expandability have allowed for the seamless integration of additional, non-identical server nodes to increase the capacity of the HPC. All controlled by a single head node, there are two independent compute environments defined by separate parallel environments within the Sun Grid Engine (SGE) job control software. Job submission and control are managed by SGE and destination queues are specified by users at the time of job submission. Within the compute nodes, the OS, working, and scratch directories are all installed on a local raid stripe while the SGE and Engineering application binaries are located physically on the cluster head node and shared with the compute nodes using NFS. The user home directories are on a separate NFS mounted share and the head node connects through fiber to a local SAN device for file storage. Infiniband interconnect hardware connects all of the cluster nodes and provides the high bandwidth backbone necessary for the simulation processing performance required to achieve the current benchmark results. A general CAE layout utilizing cluster architecture is shown in Figure 2.
2. Accelerated Simulation Performance through High Performance Computing for Advanced Sealing Applications

2.1 Advanced Sealing Applications

A Majority of sealing analysis jobs use Abaqus general-use gasket elements or elements with thickness-direction behavior only. Gaskets may be composed of materials that are softer than the materials that compose the neighboring components or it can be metal gaskets with embossments on them which perform the sealing function (Figure 3).
Discretization of gaskets will usually be finer than the discretization of mating components. This situation results in mismatched meshes being used in analyses involving gaskets. Since in most cases mismatched meshes may be the faster option, the pressure distribution may not be captured accurately and sub-modeling may be required. In most cases, two types of contact techniques are used to connect gasket elements to other parts in the model when surface-based constraints are used, namely regular contact pair or tied contact. In cases when 1D gasket elements are used, both the gasket interfaces are tied to the mating hardware component. In other cases when 3D gaskets are used, one side of the gasket is either tied using a *TIE option or contact pair that does not allow for separation. Changing this interaction behavior makes big differences in the run times needed to complete a gasket analysis for the powertrain application that can include different gaskets such as seen in Figure 4:

- Cylinder head gasket (CHG)
- Exhaust manifold gasket (EMG)
- Intake gasket
- Valve Cover gasket
- Other secondary gaskets

Figure 4: Included sealing components in a diesel powertrain
In current day powertrain applications (more predominantly in automotive industry), most of the gaskets are a highly optimized lamina of different metals and geometries often referred to as simply metal gaskets. Depending on the application and the temperature limits that the gasket sees during the operation of the engine, the metal gasket can be with or without coating. Also the specific type of coating plays a big role in the convergence of the analysis and the type of modeling techniques that can be implemented. If more detailed analysis results are required, stacked gasket elements (Figure 5) or actual 3D representation of the gasket using continuum elements (Figure 6) need to be included in the simulation.

The stacked gasket element is used when the gasket is composed of different materials or is used when relative sliding between the different layers of the gasket needs to be studied in the global model. When detailed studies are being done with a sub-model, more accurate results from the global models are being used to drive the sub-model. Here the relative slip between the different gasket layers influence how the motion on the block is transferred to the head or vice versa during, for instance a deep thermal shock test creating large temperature gradients within the powertrain. These types of detailed studies have helped to understand some of the missing information that cannot be obtained from a single element thick gasket model or even the actual dynamometer test.

Figure 5: Stacked gasket element model

3D continuum elements are used instead of gasket elements in powertrain models when the contact area cannot be represented correctly with the gasket elements. This provides a better understanding of the internal stress in the gasket material when subjected to different loads and thermal conditions that a powertrain sees. Converting a typical power train model from using gasket elements to continuum elements as shown in Figure 6, results in a large change in analysis run time due to the changes in the contact interaction area. Also, with an increase in the number of layers the analysis becomes much more complex. With the presence of soft rubber coatings on top of the metal gasket layers, there is a different coefficient of friction between a coated area and
uncoated area on the gasket. The models of this nature are used to study the internal stresses and the stress amplitudes in the gasket material and also used as an input for fatigue analysis.

![Gasket model with continuum elements](image)

**Figure 6: Gasket model with continuum elements**

Until recently most of the structural analysis was done on a component basis or joint by joint basis due the cost of computation power and the software limitations. But now with the multi-core computer technology and the DMP technology of the software, larger and more detailed models can be efficiently created and analyzed within a much shorter time frame. The faster turnaround time also helps in incorporating design-of-experiments (DOE) which in turn fine tunes design aspects not only for the gasket, but for the components and material selection as well. As more and more capabilities are incorporated into software and multi-core machines are finding their way into desktops, codes such as MpCCI can be used to couple FEA code (Abaqus) and CFD Code (Fluent) to be used for large scale FSI simulations. One application example is the creation of a thermal map of the powertrain model. Incorporating the exhaust system models into a base engine powertrain model, results in a larger model size and more complex interaction. One of the powertrain models (PT2) that is benchmarked in this paper is a very finely meshed hardware model that is used for component durability studies, which incorporates very refined meshes in certain critical locations with an element size of 1-2 mm.

### 2.2 High Performance Computing

As shown, DMP compute clusters deliver unprecedented performance for finite element applications at a relatively low cost. The approach taken to build today’s compute clusters is to use a large number of low-cost processors or cores to attack a single problem. Distributed memory architectures are central to this approach for two reasons:

- Such architectures are fundamentally scalable
• Relatively low cost because the hardware has no burden of sharing memory between a large numbers of cores.

The downside of compute clusters is that software must be specifically written to execute on distributed memory architectures which presents a major software development challenge.

A key hurdle in enabling distributed memory execution for a non-linear implicit finite element application like Abaqus/Standard is the solution of a large system of linear equations that is repeated with each solution iteration. This problem can be tackled with either direct linear equation solvers which are very robust but difficult to parallelize, or with an iterative equation solver which generally parallelize well but may not converge for many problems. With Abaqus 6.7 the SIMULIA development team released a direct linear equation solver that executes efficiently on distributed memory architectures. SIMULIA has refined this technology with Abaqus 6.8 and 6.9 and will continue to enhance this solver with future releases.

Two directions seen in today’s hardware have driven and enabled SIMULIA to approach delivering high performance for Abaqus/Standard. The first is that while hardware vendors have lost the ability to cost-effectively deliver significant speedups from scalar (non-parallel) processors, the vendors have been successful in delivering many cores in a single processor at a similar cost to the single-core processors of several years back. The second relates to networking technology, commonly referred in the HPC community as interconnect technology, having made great strides allowing large volumes of data to be moved quickly between processors. Networks continue to improve at the same time that hardware vendors continue to pack more cores into each processor which opens doors for parallel applications to deliver better performance at a fixed cost.

2.3 Performance Studies

In this paper we study the cost of running three different benchmarks on a compute cluster using a range of cores. Model images and details can be found in Figure 7 and Table 1, respectively. The compute cluster specifications that were used at R Systems NA, Inc. are presented in Table 2. All runs were done using the 6.9-PR4 release of Abaqus which is the second and final Beta release for 6.9.

![Figure 7: Benchmarks used in this study: a) Axle assembly; b) Powertrain](image)
Table 1: Pertinent benchmark information

<table>
<thead>
<tr>
<th></th>
<th>Axle</th>
<th>Powertrain 1 (PT1)</th>
<th>Powertrain 2 (PT2)</th>
</tr>
</thead>
<tbody>
<tr>
<td>Num. Elements</td>
<td>1.20E+06</td>
<td>8.75E+05</td>
<td>1.94E+06</td>
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<tr>
<td>Num. Nodes</td>
<td>1.92E+06</td>
<td>1.50E+06</td>
<td>3.12E+06</td>
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<td>Num. Equations (DoFs)</td>
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<td>5.28E+06</td>
<td>9.30E+06</td>
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<td>FLOPS/Iteration</td>
<td>3.32E+13</td>
<td>2.57E+13</td>
<td>1.60E+14</td>
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<tr>
<td>Memory to minimize I/O on one node (GB)</td>
<td>71.8</td>
<td>44.3</td>
<td>128.3</td>
</tr>
<tr>
<td>Steps</td>
<td>10</td>
<td>3</td>
<td>8</td>
</tr>
<tr>
<td>Increments</td>
<td>10</td>
<td>14</td>
<td>34</td>
</tr>
<tr>
<td>Iterations</td>
<td>65</td>
<td>47</td>
<td>154</td>
</tr>
</tbody>
</table>

Table 2: Compute cluster specifications

<p>| | |</p>
<table>
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<tbody>
<tr>
<td>Processor:</td>
<td>Intel® Xeon® Quad-core CPU Model E6462 @ 2.80GHz</td>
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<tr>
<td>Node Configuration:</td>
<td>2 socket compute nodes with 8 cores/node</td>
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<tr>
<td>Memory:</td>
<td>32GB/node</td>
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<tr>
<td>Number of Nodes:</td>
<td>32</td>
</tr>
<tr>
<td>Interconnect:</td>
<td>InfiniBand (Quad Data Rate)</td>
</tr>
<tr>
<td>Disks:</td>
<td>One 250GB SATA drive on each node</td>
</tr>
<tr>
<td>Operating System:</td>
<td>Redhat Linux Enterprise Server 5.1</td>
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</table>

All three models were run on a number of cores ranging from one to 128, however not all analyses were run to completion due to very long run times. The terminated analysis were run long enough to collect average iteration times which is the most important contributor to the total wall clock time. The average iteration times were then used to estimate the total wall clock time. Table 3 shows the total analysis wall clock times, average iteration time, and the percent contribution of iteration time for completed runs. The percent contribution of iteration time to total wall clock time for low core count runs is over 98% which indicates our total wall clock time estimates for incomplete runs is very accurate.
Table 3: Performance data for all three benchmark models

<table>
<thead>
<tr>
<th></th>
<th>Number of cores</th>
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<td></td>
<td>1</td>
<td>4</td>
<td>8</td>
<td>16</td>
<td>32</td>
<td>64</td>
<td>128</td>
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<tr>
<td>AXLE Total Wall</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>(hrs)</td>
<td>107.9</td>
<td>50.3</td>
<td>42.7</td>
<td>20.7</td>
<td>7.5</td>
<td>4.5</td>
<td>3.4</td>
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<tr>
<td>Avg. Iter.</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
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</tr>
<tr>
<td>(min)</td>
<td>98.8</td>
<td>45.6</td>
<td>38.6</td>
<td>18.0</td>
<td>6.1</td>
<td>3.6</td>
<td>2.6</td>
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<td>98%</td>
<td>94%</td>
<td>89%</td>
<td>88%</td>
<td>84%</td>
</tr>
<tr>
<td>PT1 Total Wall</td>
<td></td>
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<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>(hrs)</td>
<td>49.6</td>
<td>16.2</td>
<td>12.8</td>
<td>4.3</td>
<td>2.4</td>
<td>1.4</td>
<td>1.1</td>
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<tr>
<td>Avg. Iter.</td>
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<td></td>
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<tr>
<td>(min)</td>
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<td>20.4</td>
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<td>5.2</td>
<td>2.9</td>
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<td>–</td>
<td>99%</td>
<td>98%</td>
<td>95%</td>
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<td>90%</td>
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<tr>
<td>PT2 Total Wall</td>
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<tr>
<td>(hrs)</td>
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<td>Avg. Iter.</td>
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<tr>
<td>(min)</td>
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<td>153.4</td>
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<td>Iter. Contrib.</td>
<td>–</td>
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<td>–</td>
<td>99%</td>
<td>99%</td>
<td>98%</td>
<td>97%</td>
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</table>

Shaded cells indicate estimated total wall clock time.

Another way of looking at parallel performance is to measure the speed-up of the code when more cores are used. Figure 8 shows the speed-up of Abaqus/Standard iteration time against the number of cores used on the compute cluster. Considering the smallest amount of resources that can be allocated and charged for in a cluster environment is a single compute node, it makes more sense to measure speed-up with respect to one compute node performance (i.e., 8 cores). The super-scaling observed in Figure 8 is due to the availability of memory and parallelized I/O performance on multiple nodes. Considering the amount of memory required to minimize I/O on a single compute node (see Table 1), one can see that the PT2 model would start running totally “in-core” when 64 cores or 8 nodes are used leading to the almost 2.6X speed-up from 32 cores. The impressive scalability of PT2 simulation is particularly important considering the run time is reduced from 15 days on 8 cores to almost 11 hrs on 128 cores.
Using the timing data we are able to calculate the simulation costs. While only normalized numbers are presented in this paper, actual software license costs, analysis token requirements, and pricing for the compute cluster were used. The software license cost was based on the assumption that all tokens were used at all times since the annual license fee was used to calculate the software cost per hour. Compute cluster usage costs were more straightforward to calculate since it is simply by node hour used. Using the software and compute cluster costs per hour, an aggregate hourly cost can be calculated. Then, for each analysis the total cost can be calculated. Similarly the cost per iteration can also be calculated. Using this method, normalized iteration costs were calculated for all three benchmarks from one to 128 cores and the results are shown in Figure 9. The results show that for the AXLE and PT1 models the lowest cost of simulation is achieved at 32 cores, but for the PT2 model 64 cores is the most optimized configuration. In terms of actual savings, it is interesting to note that for PT2 the cost of running on one core is 7.7 times more expensive than running on 128 cores while the job runs 87.6 times faster.

Figure 9 shows that the cost curves for all three benchmarks have a very steep down slope at low core counts but remain relatively flat at high core counts. In simple terms this means that there is high penalty for executing jobs on low core counts but there is relatively little penalty for running on too many cores. For our benchmarks this means that the analyses can be completed almost two times faster at almost no extra cost compared to the sweet spot (i.e., optimal case).
Figure 9: Relative costs per iteration

2.4 Updated simulation structure

In an ideal form, parallel hardware and software fundamentally change the choices finite element analysts have to make about model size and runtime. Previously, when executing on a single processor host the only choice available to an analyst wanting to run a larger model or run an existing model faster, was to wait for the next generation of hardware that would deliver better speeds and larger memory sizes. For an analyst executing perfectly scalable software on perfectly scalable hardware, model sizes and execution speeds are determined purely by how many individual compute hosts the analyst is willing to dedicate to the job rather than by the latest revision of the hardware. In reality the equation is more complicated because Abaqus is not perfectly scalable. As hardware and software improve, Abaqus scaling will improve. However, even today, many analysts running large models (more than 5 million degrees of freedom) are not extending their computation allocations beyond the range of the Abaqus scaling previously shown.

Analysts should understand the power given to them by high performance computing. Consider the case of a user who typically runs Abaqus jobs on 16 cores with a turnaround of 24 hours, but one day ends up with a need for fast turnaround to handle a time-critical workload. If the compute power is available the analyst could choose to run the job on 64 cores and get a turnaround of 8 hours. For many analysts this flexibility might prove valuable, but the hardware necessary to allow this flexibility may be problematic. Utility computing providers like R Systems Inc provide a solution to this problem. A company might choose to have enough hardware in house to handle
routine compute needs, but choose to relieve a computing peak by moving computation to an off-site resource provider.

One obstacle to doing computations at remote locations still remains - post processing the model. A large Abaqus analysis may generate tens of gigabytes of output data. Transferring this data from a remote compute site back to the user’s local workstation over the Internet may take 12-24 hours which often negates the benefits obtained by making use of remote computing resources. A solution to this problem is emerging in a variety of products, both commercially available and open sourced. The solution is to do the graphics processing on a machine co-located near the remote compute cluster and then send compressed images back to the user’s machine. Although for Abaqus post processing solutions tailored to 3D graphics must be used, the approach is very similar to the commonly used Windows Remote Desktop or VNC desktop solutions.

In practice there is a minimum network bandwidth that must be available to employ these types of remote graphics solutions. A basic rule of thumb is that 1 megabit per second of bandwidth is required for every frame per second a user wants. SIMULIA engineers have done some initial testing with this approach to post processing and have found that Abaqus/Viewer can be reasonably used for large models with bandwidths of about 3-5 Mbits/sec.

The advances in compute system capabilities have quickly brought on an increased need for remote visualization tools. The ideal consolidation of large simulation systems which would allow for increased computational capacity and reduced administration costs could be further accelerated through improvements in remote visualization components. Whether reviewing inputs, checking on simulations in progress, or needing to quickly view results, the need for such utilities must not be overlooked.

Today’s tools for remote visualization provide an efficient way for engineers to interact graphically with their simulations at locations other than a directly connected terminal monitor. The network ODB viewer which is built into Abaqus and third-party tools such as Hewlett-Packard’s Remote Graphics software provide the basic requirements for the remote visualization of Abaqus simulations. Still, further developments of remote visualization tools could provide geographically dispersed system users to share a larger compute system with a perceived experience of being located in front of the actual system. Graphically, object manipulation is a limiting factor of the current tools used for remote visualization. Influences such as network bandwidth and system resources on both the client and server sides of the visualization tools have a large effect on the overall remote visualization experience. Not surprisingly, the further both geographically and by total number of interim network “hops” away from the data source that a remote user is located, the more degraded the visualization experience will be. With acceptable usability bandwidth requirement of between 3 and 5 Mbits/sec for today’s remote visualization tools, the supporting environment is limited mainly to local LAN setups. In a business environment where network resources are in competition with other business process, an ideal solution is to have a separate LAN dedicated to this purpose.
2.5 Future technology

Over the past 5 years SIMULIA has been working to re-architect the Abaqus products to run on distributed memory architectures after many years of single processors and then shared-memory parallel execution. Users may ask whether a new architecture is going to emerge in the next 5 to 10 years. The answer to this question is yes and no. Parallel computing continues to be the most obvious avenue to delivering unlimited performance, and both software developers and hardware vendors continue to learn how to better deliver parallel solutions to end users. Improvements in the performance of individual scalar compute cores continue to be made, but at a very slow rate.

A major concern in parallel computing is power consumption, and this is one of the focus areas of future technology. One approach is to use lower-speed, lower-power processors, but to use more of them as the power savings to be gained from lowering clock speeds more than offsets the use of more cores. This approach is represented by IBM’s BlueGene supercomputer and by the products of companies like SiCortex which couple relatively low speed processors with fast interconnects to deliver very good performance for applications that can effectively exploit a highly parallel architecture.

A second approach to lowering power consumption and delivering greater performance is to use graphics processors for general purpose computing (GPGPU). This represents a significant software challenge, but the performance gains that can be delivered by executing appropriate code on GPGPU’s are significant. Currently codes being tested on GPGPU’s typically execute using a single GPGPU (the code executing on the GPGPU runs many parallel streams of execution), but going forward it is an obvious step to start deploying many GPGPU’s to work in parallel on single problem.

3. Conclusion

In order to meet today’s fast-paced development requirements, up-front simulation techniques play a major role. As such, for an efficient CAE implementation to be successful, the right system environment is needed. The optimal performance must be known for the specific product portfolio as the diversity of models have significant influence on the sweet spot. On top of that the performance of the software for different hardware structures varies.

In order to meet the challenges of our time for advanced sealing solutions in the automotive and heavy duty environment, the synergy between the Abaqus FEA solver and multi-node HPC clusters is used to minimize the unit cost and retain the optimal turn-around time. Exceptional is the fact that with Abaqus V6.9 the costs are not increasing with respect to the future powertrain model (PT2) having a sweet spot of 64 cores. This is important when scheduling the work load based on timing needs in order to decrease turn-around time. In the past after reaching the sweet spot, costs would increase again. In our case the cost are almost flat when reaching 32 cores making the life of an analyst, who is always fighting to meet deadlines, much easier.

Further optimization of the cost structure and performance can be reached through remote visualization. The technology and infrastructure though have still to be improved in order to make this a viable solution in a cost effective way.
4. References